

## Solution Final Spring 2016

Please write your answers on these sheets, use the back sides if needed. Show your work. You can use a fact from the slides/book without having to prove it unless you are specifically asked to do so. Be organized and use readable handwriting. There is a page for scratch work at the end.

**Exercise 1 (Cholesky decomposition.)** (10 pts) Suppose that the square matrix  $B$  has the QR-factorization  $B = QR$ , where  $Q$  is orthogonal and  $R$  is upper triangular with positive diagonal terms, and another matrix  $A = B^\top B$ . Describe the Cholesky method for solving the system of equations  $Ax = b$  for some vector  $b$  using this information.

$A = B^\top B = R^\top Q^\top QR = R^\top R$ . Thus  $Ax = R^\top Rx = b$ . Let  $z = Rx$ . First, solve  $R^\top z = b$ , which is quick as  $R$  is triangular. This gives  $z$ . Second, solve  $Rx = z$ , which gives  $x$ . This is again quick as  $R$  is triangular.

**Exercise 2 (Ridge regression.)** (10 pts) Ridge regression involves solving optimization problems of the form  $\min \|Ax - b\|_2^2 + \lambda \|x\|_2^2$ , where  $\lambda$  is a positive regularization parameter and  $A$  and  $b$  are given. Write this problem as an equivalent least-squares problem (without regularization) such that standard least-squares methods can be used.

It is clear that  $\|Ax - b\|_2^2 + \lambda \|x\|_2^2 = \|Cx - d\|_2^2$ , where  $C = (A; \sqrt{\lambda}I)$  and  $d = (b; 0)$ .

**Exercise 3 (Risk minimization.)** Let  $c_j \in \mathbb{R}^n$  be a vector of costs faced in a future scenario  $j$ , with  $j = 1, \dots, m$ , and  $x \in \mathbb{R}^n$  be a decision vector to be optimized subject to the constraints  $x \geq 0$  and  $\sum_{i=1}^n x_i = 1$ . Since we are unsure about the future scenario, we adopt a risk-based formulation where we aim to minimize the  $\alpha$ -superquantile. This leads to the problem

$$\min f(\bar{x}) = \xi + \frac{1}{1 - \alpha} \frac{1}{m} \sum_{j=1}^m \max\{0, c_j^\top x - \xi\}, \text{ subject to the constraints,}$$

where  $\bar{x} = (\xi, x) \in \mathbb{R}^{n+1}$ . This problem can be reformulated as a linear program with  $n+m+1$  variables,  $2m+n$  inequality constraints and one equality constraint. There are algorithms that can solve this linear program in computational effort that is proportional to  $(m+n)^{3.5}$ . In this question we contrast this effort with that of the subgradient method.

The subgradient method for this problem is as follows. Starting from an initial point  $\bar{x}_0 \in \mathbb{R}^{n+1}$ , the function  $f$  is minimized directly without reformulation using the recursion  $\bar{x}_{k+1} = P(\bar{x}_k - s_k \nabla f(\bar{x}_k))$ ,  $k = 0, 1, 2, \dots$ , where  $\nabla f(\bar{x}_k)$  is a subgradient of  $f$  at  $\bar{x}_k$ ,  $s_k$  is a step size, and  $P(\cdot)$  is the projection onto the feasible set.

1. (10 pts) Since one can use a chain rule to obtain a subgradient of  $f$  at  $\bar{x}_k$ , it is sufficient to work out a formula for the subgradient of a function  $g_j(\bar{x}) = \max\{0, c_j^\top \bar{x} - \xi\}$  at  $\bar{x}_k = (\xi_k, x_k)$ . Write such a formula.

If  $c_j^\top x_k - \xi_k > 0$ , then a subgradient is  $(-1, c_j^\top)^\top$ , otherwise it is zero.

2. (5 pts) Suppose  $n$  is relatively small, but  $m$  is huge. Would you prefer the linear programming approach described above or the subgradient method for solving the above problem. Explain why.

For large  $m$ , the LP approach will be extremely costly. The subgradient approach might take many iterations, but each iteration is very cheap, roughly of order  $m$ . I would prefer the subgradient method (in reality accelerated versions of this algorithm).

**Exercise 4 (Steepest descent method.)** (10 pts) It is known that iterations of the steepest descent method generate a sequence of iterates  $x_k$  that are gradually closer to optimality with progress bounded by the expression

$$f(x_{k+1}) - p^* \leq c(f(x_k) - p^*),$$

where  $f$  is the (convex) function being minimized,  $p^*$  its minimum value, and  $c \in (0, 1)$  is a constant. Derive a formula in terms of the initial error  $\epsilon_0 = f(x_0) - p^*$ , with  $x_0$  being the initial solution, for how many iterations it takes this method to reach an  $x$  with  $f(x) - p^* \leq \epsilon$ , where  $\epsilon > 0$ .

After  $k$  iterations, we by recursion that  $f(x_k) - p^* \leq c^k(f(x_0) - p^*)$ . Thus, we needed  $c^k(f(x_0) - p^*) \leq \epsilon$  or

$$k \geq (\log c)^{-1} \log \frac{\epsilon}{\epsilon_0}$$

**Exercise 5 (Barrier method.)** We consider the original problem  $\min f_0(x)$  subject to  $f_i(x) \leq 0$ ,  $i = 1, \dots, m$ , and the corresponding logarithmic barrier problem  $\min f_0(x) - (1/t) \sum_{i=1}^m \log(-f_i(x))$ , where  $t > 0$ .

1. (10 pts) Write the KKT conditions for the original problem.

Primal feasibility:  $f_i(x) \leq 0$  for all  $i$

Dual feasibility:  $\lambda_i \geq 0$  for all  $i$

Complementary slackness:  $\lambda_i f_i(x) = 0$  for all  $i$

Gradient:  $\nabla f_0(x) + \sum_i \lambda_i \nabla f_i(x) = 0$

2. (5 pts) Write the KKT conditions for the barrier problem.

$$\nabla f_0(x) + \sum_i \frac{1}{-t f_i(x)} \nabla f_i(x) = 0.$$

3. (5 pts) Given  $t > 0$ , let  $x^*(t)$  be an optimal solution of the barrier problem and set  $\lambda_i^*(t) = -1/(t f_i(x^*(t)))$ ,  $i = 1, \dots, m$ . Suppose that  $f_0, f_1, \dots, f_m$  are convex. Show that  $x^*(t)$  is a minimum solution of the Lagrangian of the original problem when the multipliers  $\lambda_i$  are set to  $\lambda_i^*(t)$  for all  $i = 1, \dots, m$ .

Since  $x^*(t)$  is an optimal solution of the barrier problem,

$$\nabla f_0(x^*(t)) + \sum_i \lambda_i^*(t) \nabla f_i(x^*(t)) = 0.$$

In view of convexity of the Lagrangian, this implies the result.

**Exercise 6 (Support vector machine.)** (10 pts) A Support Vector Machine approach to classification needs to consider the problem

$$\min_w \frac{1}{m} \sum_{i=1}^m \hat{E}(y_i \phi(x_i)^\top w) = \frac{1}{m} \sum_{i=1}^m \max\{0, 1 - y_i \phi(x_i)^\top w\}$$

where  $\{x_i, y_i\}_{i=1}^m$  is given data and  $\hat{E}(\alpha) = \max\{0, 1 - \alpha\}$  is the hinge loss. Write this problem as a linear program.

$$\begin{aligned} \min_{w, e} \quad & \frac{1}{m} \sum_{i=1}^m e_i \\ \text{subject to} \quad & 1 - y_i \phi(x_i)^\top w \leq e_i, \quad i = 1, \dots, m \\ & e_i \geq 0, \quad i = 1, \dots, m \end{aligned}$$

**Exercise 7 (Shape-Constrained Regression.)** Suppose that we have some data  $\{x_i, y_i\}_{i=1}^m$ , with  $x_i, y_i \in \mathbb{R}$ . We would like to carry out least-squares regression on this data set, but

need to ensure that the regression function is convex. We will achieve this using epi-splines. First, we discretize a sufficiently large part of the first axis by the points  $m_0, m_1, \dots, m_N$ . Second, on each segment  $(m_{k-1}, m_k)$ , we let the regression function be a second-order polynomial of the form  $a_0^k + a_1^k x + a_2^k x^2$ . The  $a$  coefficients are to be determined by the optimization. Consequently, the regression function  $f_a(x)$  is a one-dimensional function defined on  $[m_0, m_N]$ , which is piecewise polynomial with coefficients determined by the various  $a$ -coefficients. Least-squares regression aims to find such a regression function that has a small error relative to the observed data.

1. (5 pts.) Write an objective function that expresses the least-squares criterion in this case. You can assume that the regression function is continuous.

Let  $e_i = y_i - [a_0^{k_i} + a_1^{k_i} x_i + a_2^{k_i} x_i^2]$ , where  $k_i$  is such that  $x_i \in [m_{k_i-1}, m_{k_i}]$ . Then the objective becomes  $\sum_{i=1}^m e_i^2$ .

2. (10 pts.) Write a set of constraints that ensures that the regression function is convex. Make sure you include constraints that enforce continuity of the regression function.

Continuity:  $a_0^k + a_1^k m_k + a_2^k m_k^2 = a_0^{k+1} + a_1^{k+1} m_k + a_2^{k+1} m_k^2$  for  $k = 1, \dots, N - 1$ .

Convexity at mesh points:  $a_1^k + 2a_2^k m_k \leq a_1^{k+1} + 2a_2^{k+1} m_k$  for  $k = 1, \dots, N - 1$ .

Convexity in segments:  $a_2^k \geq 0$  for all  $k = 1, \dots, N$ .

**Exercise 8 (Control.)** (5 pts.) We need to develop a control algorithm for a robot that moves at constant speed  $v$  in a two-dimensional space. The robot's motion is modeled as a Dubin's vehicle, i.e.,  $\dot{x}_1(t) = v \cos x_3(t)$ ,  $\dot{x}_2(t) = v \sin x_3(t)$ , and  $\dot{x}_3(t) = u(t)$ , where  $x_1(t)$  and  $x_2(t)$  are the coordinates in the plane at time  $t$  and  $x_3(t)$  is the heading at time  $t$ . The control input at time  $t$  is  $u(t)$ . Consider Euler's method of solution of this differential equation with time discretization step  $\Delta t$  so that the discretized version of the optimal control problem will involve the state variables  $x_i(k\Delta t)$ , for  $i = 1, 2, 3$  and  $k = 0, 1, 2, \dots, N$ .

Using the state variables at the discretized points in time, write a set of constraints that ensures that the robot is never further away from a desired trajectory given by  $\{(y_1(t), y_2(t)), t \geq 0\}$  than the robot's second coordinate at the same points in time. What type of constraints will this be?

Let  $x(t) = (x_1(t), x_2(t))$  and  $y(t) = (y_1(t), y_2(t))$ . Then, the constraints will be  $\|x(k\Delta t) - y(k\Delta t)\|_2 \leq x_2(k\Delta t)$  for  $k = 0, 1, \dots, N$ , which are second-order cone constraints.

**Exercise 9 (Positive definiteness.)** (5 pts) Consider an  $n$ -by- $n$  symmetric matrix with smallest eigenvalue  $\lambda_{\min} > n$ . Suppose that this matrix is augmented with one row at the

bottom and one column to the right consisting exclusively of ones. The augmented matrix is then of dimension  $(n + 1)$ -by- $(n + 1)$ . Recall that an  $m$ -by- $m$  matrix of only ones has largest eigenvalue of  $m$ . Prove that the augmented matrix is positive definite.

The Schur complement theorem establishes that it is sufficient to prove that  $A - B$  is positive definite, where  $A$  is the original  $n$ -by- $n$  matrix and  $B$  is an  $n$ -by- $n$  matrix of ones. Since  $y^\top(A - B)y = y^\top Ay - y^\top By \geq \lambda_{\min}\|y\|_2^2 - n\|y\|_2^2 > 0$  by the Rayleigh quotient theorem for nonzero  $y$ , we have established to conclusion.